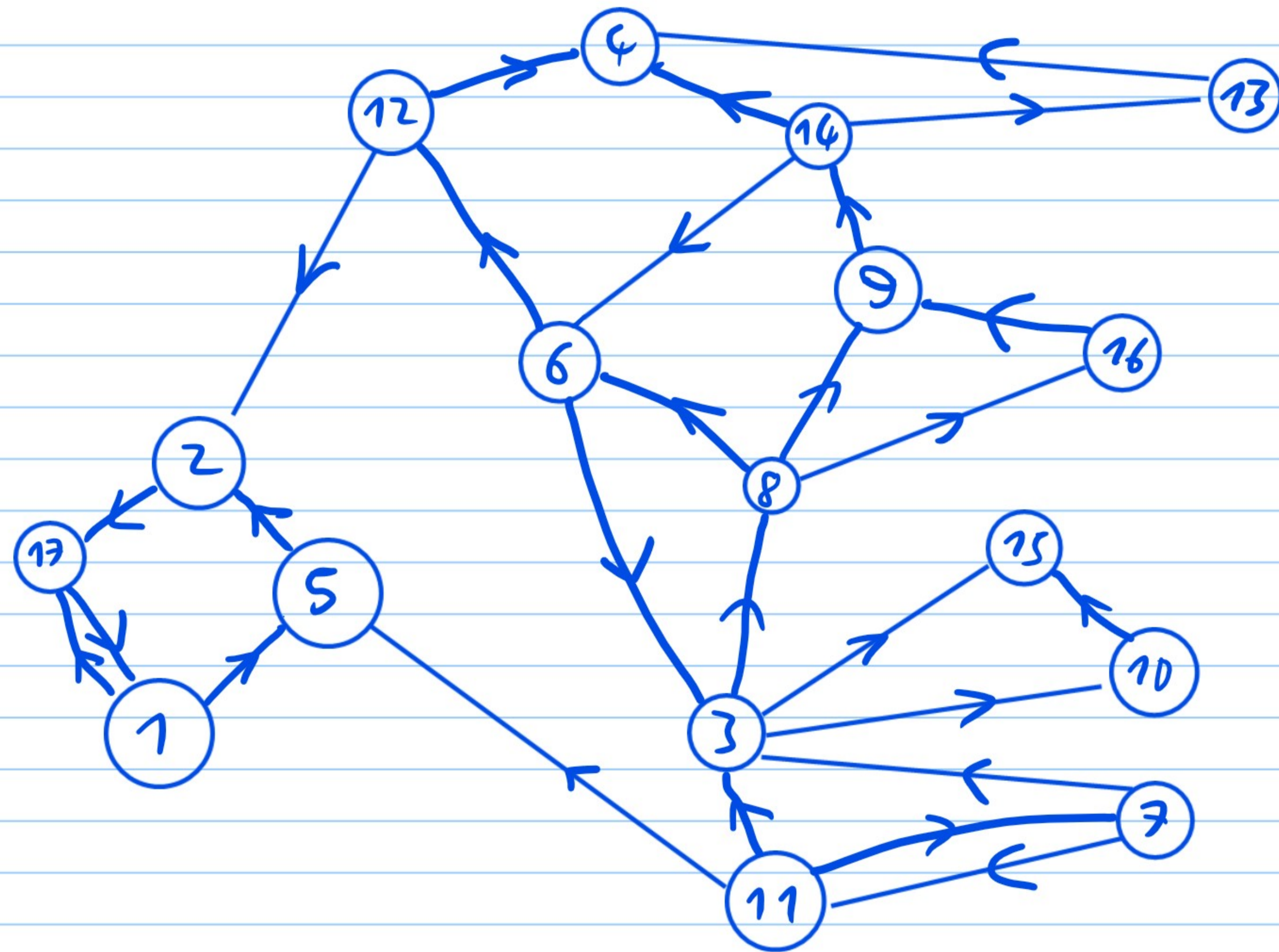


# The scaling limit of a critical random directed graph

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## I Introduction

1)  $\vec{G}(n, p)$



vertices:  $\{1, 2, \dots, n\}$

edges: keep each of the  $n(n-1)$  possible  
directed edges independently with  
probability  $p$ .



We're interested in **strongly connected components**  
(SCC)

$x$  and  $y$  are in the same SCC if  
there are two paths:  $x \rightarrow y$   
 $y \rightarrow x$

## b) Phase transition for the SCC

Same as the Erdős-Rényi graph:

$p = \frac{c}{n}$

- $c > 1 \rightarrow$  unique giant component
- $c < 1 \rightarrow$  only small components

In fact, something more precise is known

Theorem: Luczak & Seierstad (09)

Write  $p = \frac{1}{n} + \frac{d_n}{n^{4/3}}$ , assume  $d_n = o(n^{2/3})$

- If  $d_n \rightarrow +\infty$  then:
  - largest SCC has size  $\sim 4n^{2/3} d_n^2$
  - 2nd largest has size  $O(n^{2/3}, d_n)$
- If  $d_n \rightarrow -\infty$  then:
  - largest SCC has size  $O\left(\frac{n^{2/3}}{|d_n|}\right)$



Q: What happens if  $d_n$  stays bounded or converges?

II Case of the Erdős-Rényi graph

$$G(n, p) \quad p = \frac{1}{n} + \frac{d}{n^{4/3}} + o\left(\frac{1}{n^{4/3}}\right) \quad d \in \mathbb{R}$$

Theorems:

• Aldous (97): let  $Z_1^n \geq Z_2^n \geq \dots$

be the sizes of the connected components of  $G(n, p)$ . Then

$$\frac{Z_i^n}{n^{2/3}} \xrightarrow{(d)} \sigma_i$$

(in  $\ell^2$ , as sequences)

• Addario-Berry, Broutin, Goldschmidt (12)

Let  $(C_i^n)_i$  be the cc of  $G(n, p)$  with



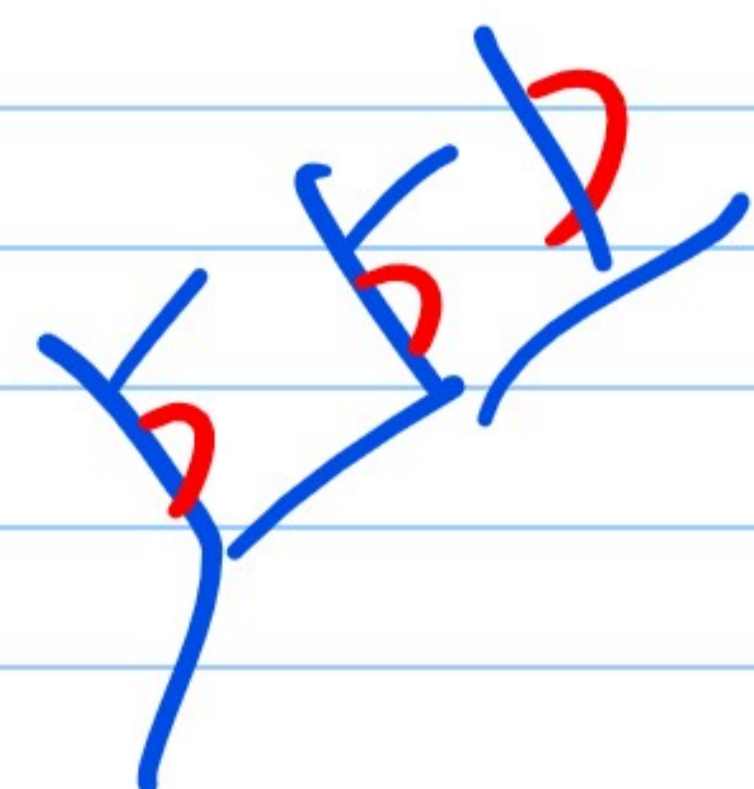
$\# C_i^n = Z_i^n$ . Then

$$\frac{C_i^n}{n^{2/3}} \xrightarrow{L_1} \mathcal{C}_i$$

(in Gromov-Hausdorff, as a sequence)

The main takeaway is that the distance between two typical points is of order  $n^{2/3}$ .

(Side note: both the  $C_i^n$  and  $\mathcal{C}_i$  have descriptions as "binary trees with a few additional edges")





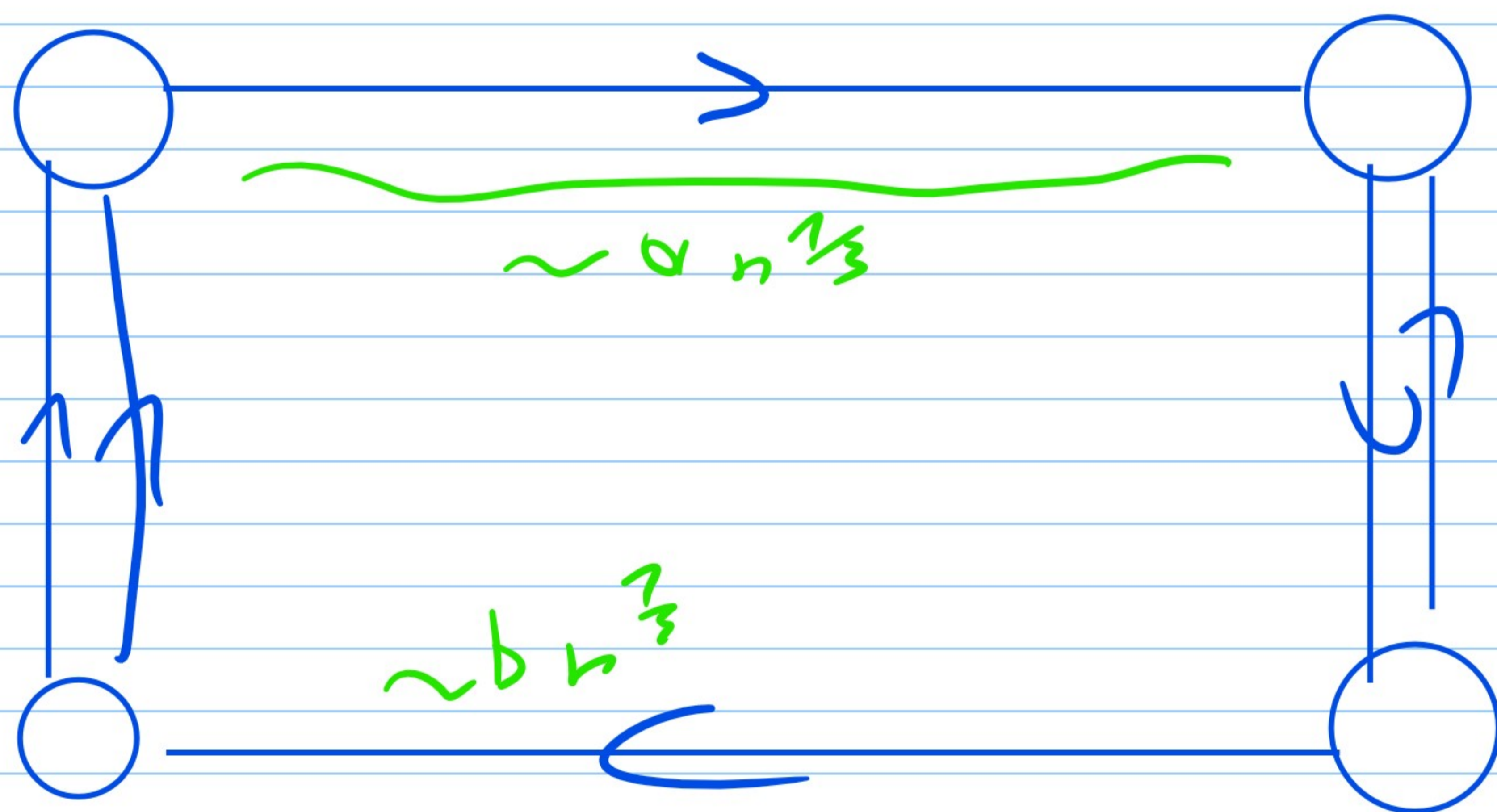
### III Our results

$$p = \frac{2}{n} + \frac{2}{n^{4/3}} + o\left(\frac{1}{n^{4/3}}\right)$$

#### a) Informally

Within the SCC:

- w.h.p. **no** vertices with degree  $\geq 4$
- number of vertices with degree=3 is  $\approx 1$
- number of vertices with degree=2 is  $\approx n^{2/3}$



#### b) Metric directed multigraphs (MDM)

definition:  $\mathcal{G}^{\rightarrow}$  is the set of finite  
multigraphs where each edge



has a direction and a length

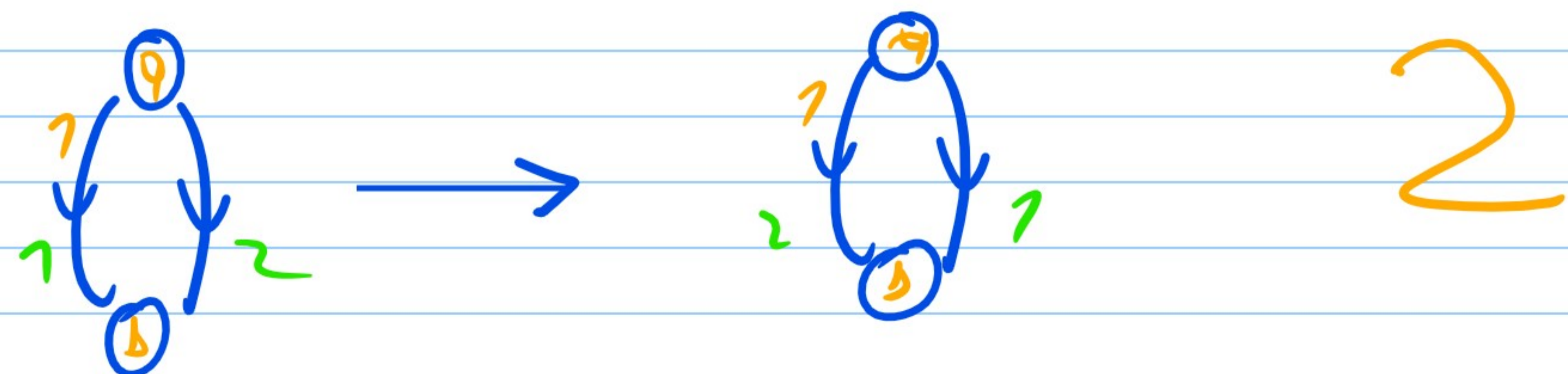
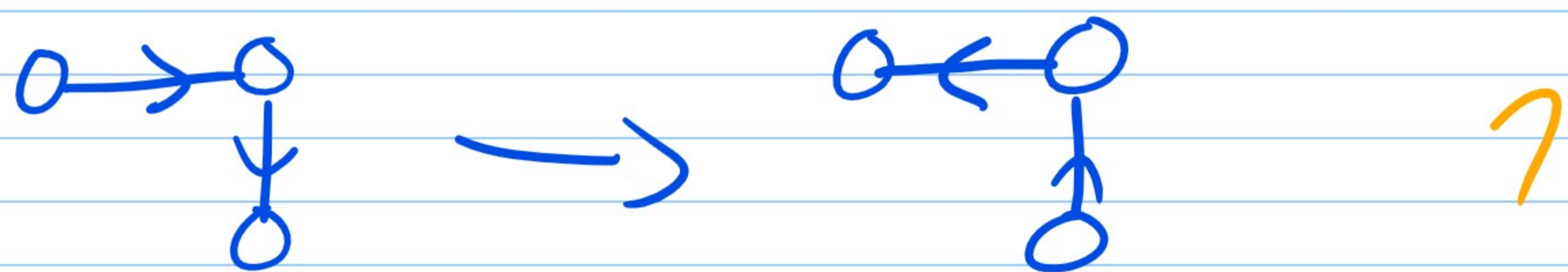
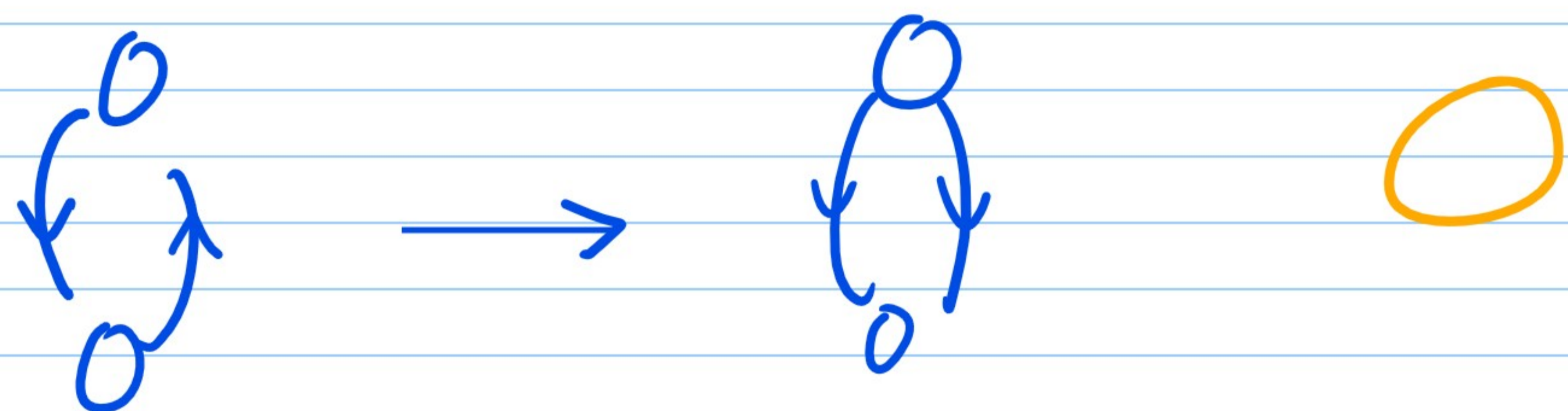
• An isomorphism is a pair of bijections

$$f: V(X) \rightarrow V(Y)$$

$$g: E(X) \rightarrow E(Y)$$

which preserve the structure

examples: How many isomorphisms?



$$\text{distance: } d_g(X, Y) = \inf_{\substack{(f, g) \\ \text{isom } X \rightarrow Y}} \sum_{e \in E(X)} |l_x(e) - l_y(g(e))|$$



$$d_{\mathcal{G}}^0 \left( \left( \begin{array}{c} \circ \\ \nearrow \\ \circ \\ \searrow \\ \circ \end{array} \right)^{2+\epsilon}, \left( \begin{array}{c} \circ \\ \nearrow \\ \circ \\ \searrow \\ \circ \end{array} \right)^1 \right) = \epsilon + \epsilon'$$

remark: This metric is very **rigid** :

$$d_{\mathcal{G}}^0 \left( \left( \begin{array}{c} \circ \\ \uparrow \\ \circ \\ \searrow \\ \circ \end{array} \right)^{\epsilon}, \left( \begin{array}{c} \circ \\ \searrow \\ \circ \\ \searrow \\ \circ \end{array} \right)^1 \right) = \infty$$

c) Main Theorem:

Let  $(C_i(n_i), n_i \in \mathbb{N})$  be the SCC of  $\vec{G}(n, p)$ , viewed as MDMs by removing the  $d^0 = 2$  vertices. Then:

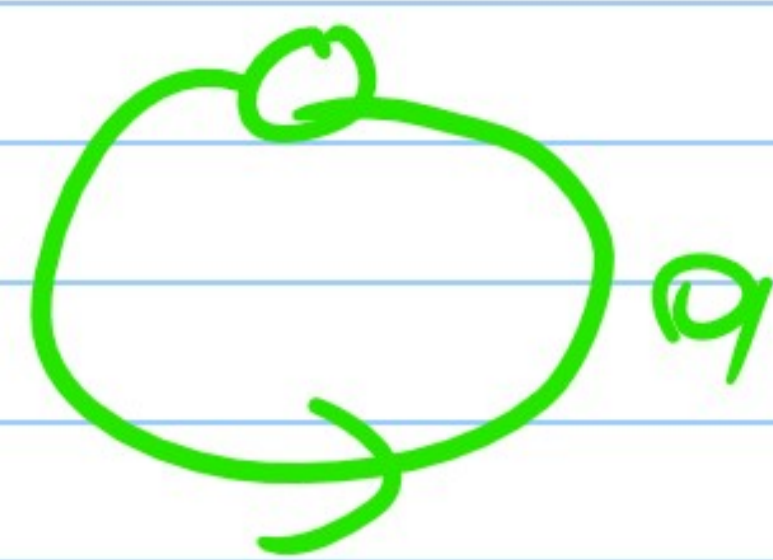
$$\left( \frac{C_i(n_i)}{n^{2/3}}, i \in \mathbb{N} \right) \xrightarrow{d} (\mathcal{C}_i, i \in \mathbb{N})$$

where  $\mathcal{C} \in \mathcal{G}^{\rightarrow \mathbb{N}}$  s.t.



• finitely many terms are 3-regular

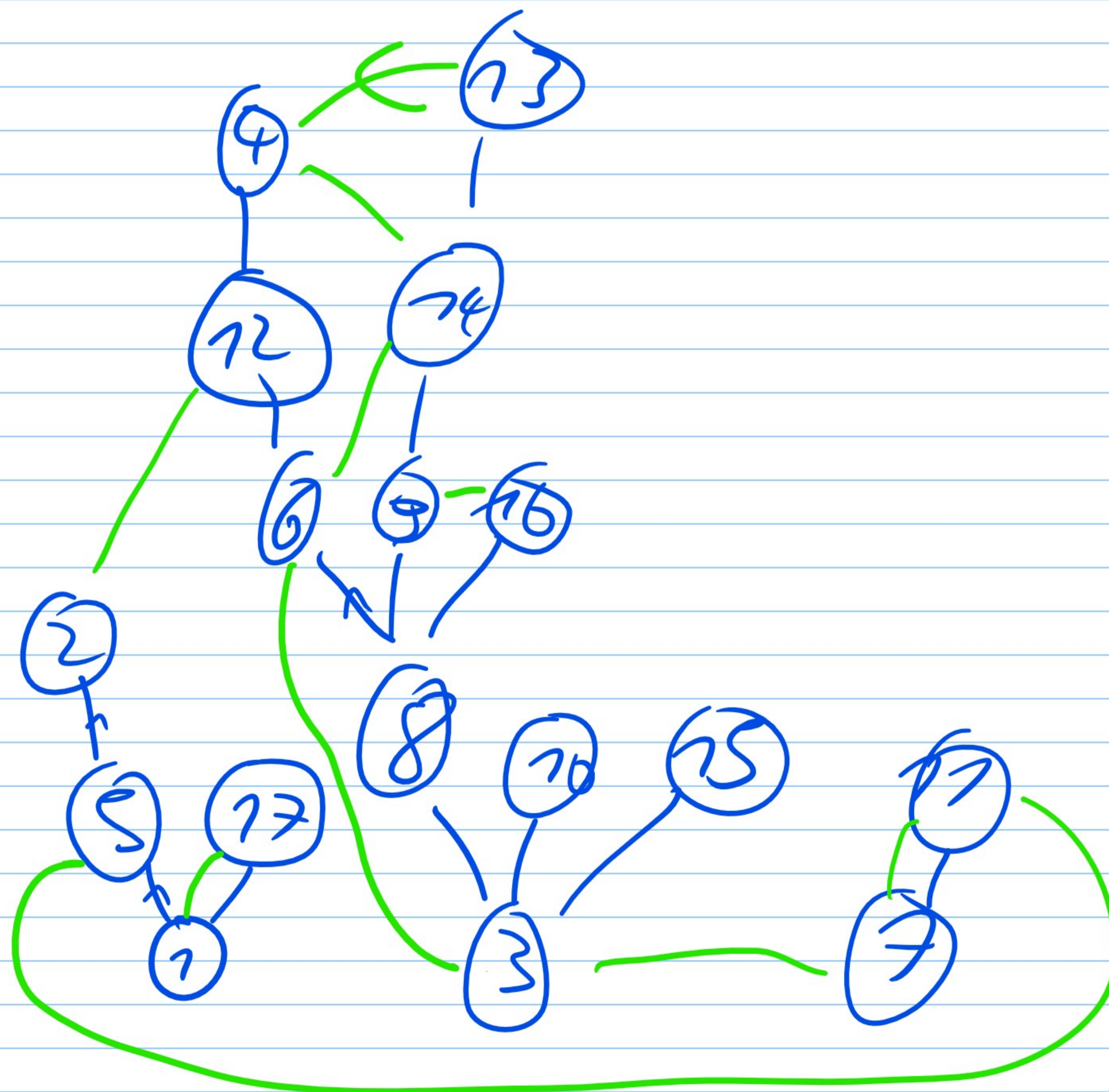
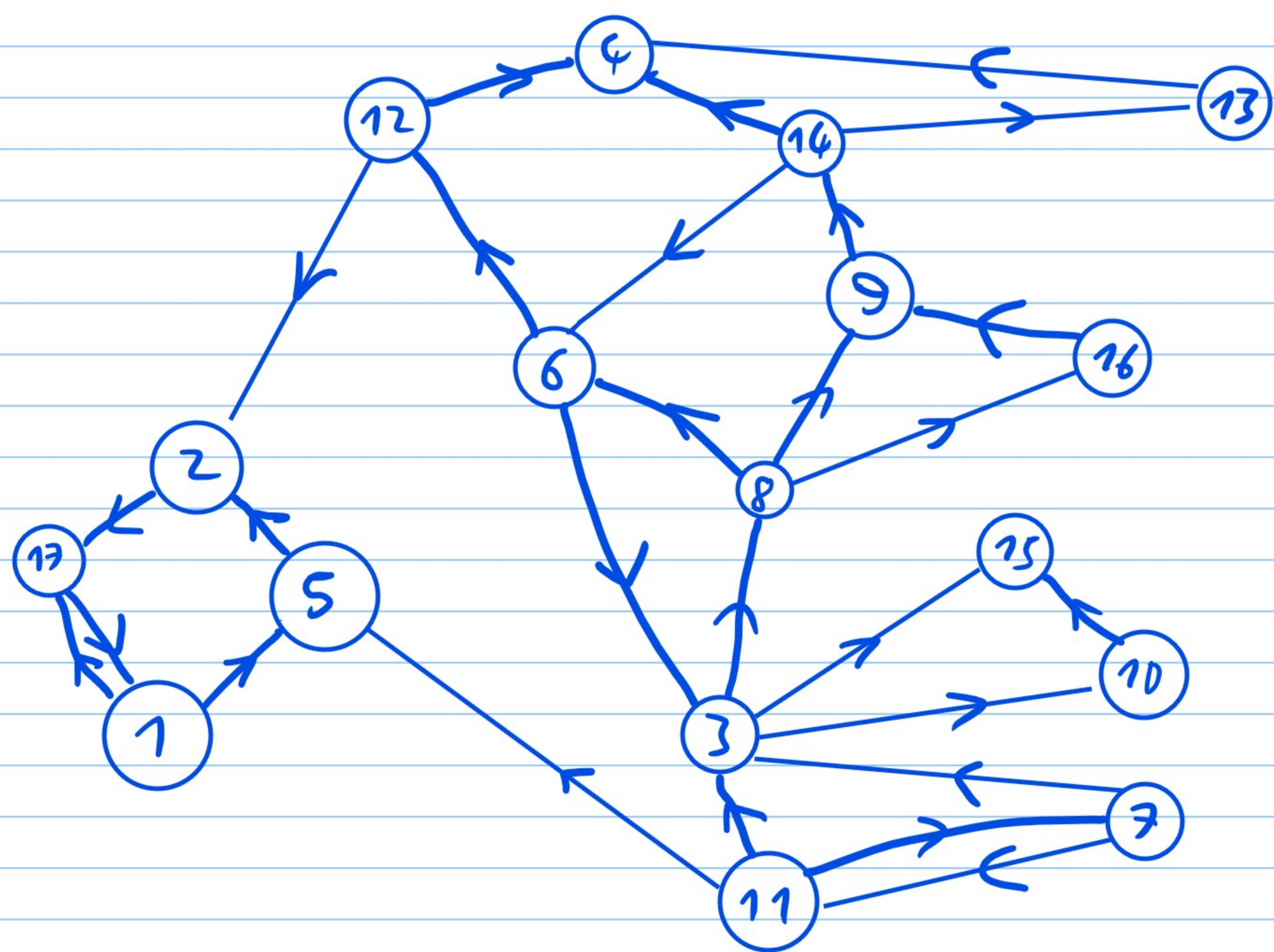
• the rest are cycles.



Topology for sequences:  $d(A, B) = \sum_{i=1}^{\infty} \frac{d_{\infty}(A_i, B_i)}{2^i}$

### IV Exploration and structure

1) A version of depth-first search





## b) Edge classification

- tree edge (forward)
- surplus edge (forward)
- back edge (backward)

Interaction between forward and backward  
→ SCC.

## ∇ Limit behaviour of forest and surplus

prop.: "Forward edges without arrows"

$$\langle d \rangle = O(n, p)$$

↗ Erdős-Rényi

So we know a lot about the trees and surplus edges. In particular:

- number of vertices in a tree  $\sim n^{2/3}$
- typical distances are  $\sim n^{1/3}$



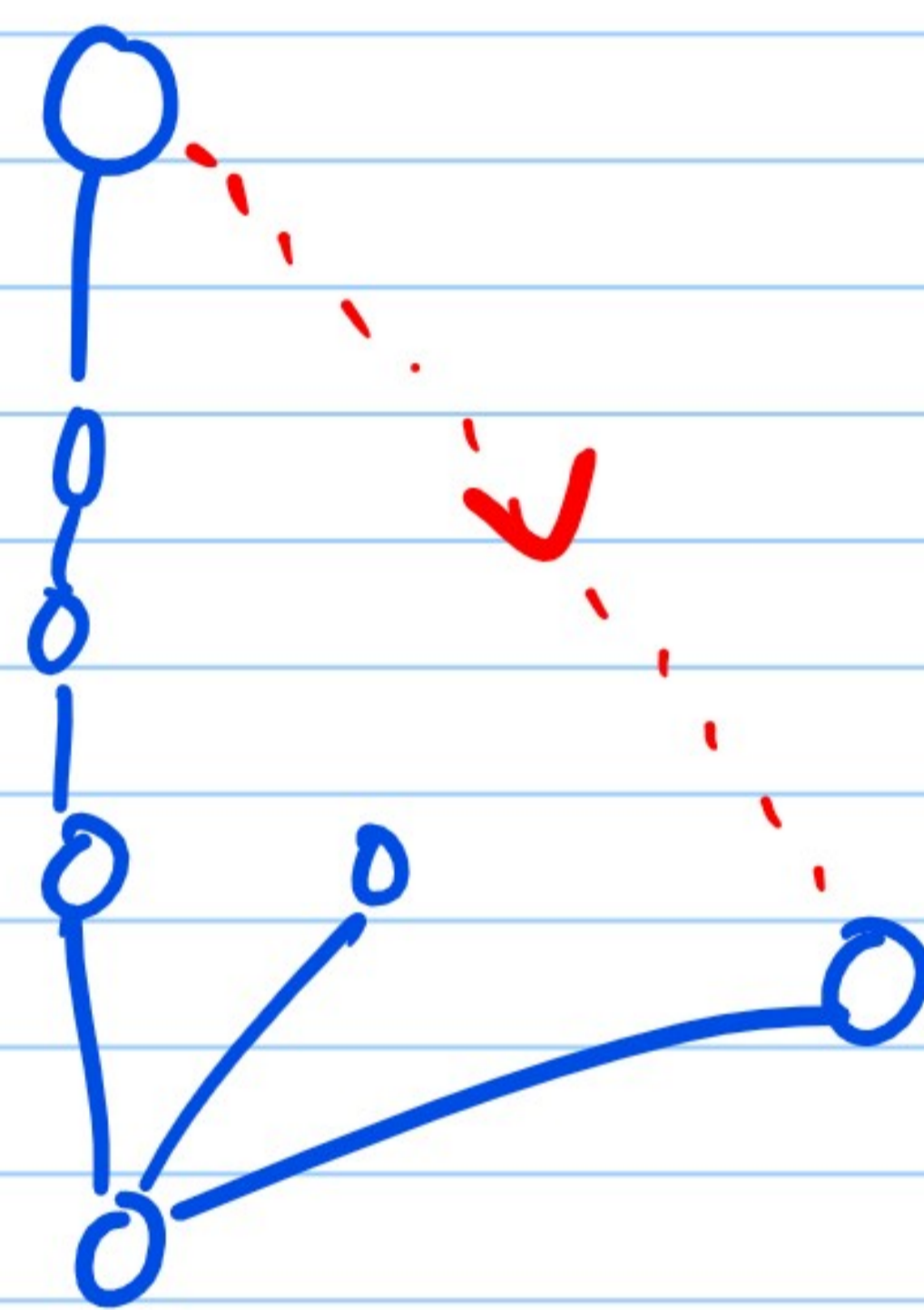
- surplus edges are  $O(1)$  in number.

It turns out, surplus edges don't count!

prop:  $P(\text{a SCC of } \vec{G}(n,p) \text{ has a surplus edge}) \xrightarrow{n \rightarrow \infty} 0$

"pf":

$P(\dots \text{ counts})$



## VI Limit of the back edges

a) problem?

In a single tree, there are  $\frac{k(k-1)}{2}$  possible back edges, appearing independently with probability  $p \approx \frac{1}{n}$ .

But  $k \approx n^{2/3} \dots$



$$E(\# \text{ of back edges}) \approx \frac{1}{n} n^{4/3} \rightarrow \infty \quad ??$$

b) solution! Most of the BE don't matter!

How to find those which do matter:

- go around the tree
- keep the first **ancestral** BE
- keep any BE which is either **ancestral** or **links into something we've already selected**.

This turns out to converge to a PPP on the limiting continuum tree!



What we end up with:



